Mad families and the modal logic of \mathbb{N}^*

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Mathematical Logic

The modal logic of $\beta(\mathbb{N})$

Guram Bezhanishvili · John Harding

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Mathematical Logic

The modal logic of $\beta(\mathbb{N})$

Guram Bezhanishvili \cdot John Harding or Rather of $\beta \mathbb{N} \setminus \mathbb{N}$

is S4 if $\mathfrak{a} = \mathfrak{c}$ or if ...

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We recall that **S4** is the least set of formulas containing the Boolean tautologies, the axioms:

 $\begin{array}{c} \Box(p \rightarrow q) \rightarrow (\Box p \rightarrow \Box q) \\ \Box p \rightarrow p \\ \Box \Box p \rightarrow \Box p \end{array}$

and closed under Modus Ponens $(\varphi, \varphi \rightarrow \psi/\psi)$ and Necessitation $(\varphi/\Box \varphi)$. Rela-

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McKinsey and Tarski defined a valuation ν of formulas of \mathcal{L} into (W, τ) by putting

- ν(p) ⊆ W,
 ν(¬φ) = W − ν(φ),
- $\nu(\neg\varphi) = w \nu(\varphi),$
- $\nu(\varphi \lor \psi) = \nu(\varphi) \cup \nu(\psi)$,
- $\nu(\varphi \land \psi) = \nu(\varphi) \cap \nu(\psi)$,
- $\nu(\varphi \rightarrow \psi) = (W \nu(\varphi)) \cup \nu(\psi)$,
- $\nu(\Box \varphi) = Int(\nu(\varphi)),$
- $\nu(\Diamond \varphi) = \overline{\nu(\varphi)}.$

In definitions and arguments in this paper, we will often economize, and leave out the clauses for disjunction, implication and modal diamond, as these are automatic from the others. Now, call a triple $M = \langle W, \tau, \nu \rangle$ a topological model. A formula φ is said to be true in such a model M if $\nu(\varphi) = W$, and we say that φ is topologically valid if it is true in every topological model. Referring to the second axiomatization of **S4**, which highlights the interior operator, one easily sees its soundness:

If $S4 \vdash \varphi$, then φ is topologically valid.

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Question

Does \mathbb{N}^* map onto every finite topological space by an open continuous map (with crowded fibers)

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[BH] proved that it suffices to work with finite T_0 -spaces providing we have crowded fibers.

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related to Veksler Problem: Can \mathbb{N}^* have maximal nwd sets?



and $f : \mathbb{N}^* \to T$

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Then (not previously known to exist in ZFC)

• U_{t0}, U_{t1}, U_{t2} are disjoint (regular) open sets



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Beszhanishvili-Harding used $\mathfrak{a} = \mathfrak{c}$, i.e. madf's

now for adf's

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Say that an open $U \subset \mathbb{N}^*$ is an adf* if there is an infinite adf \mathcal{A} such that $U = U_{\mathcal{A}} = \bigcup_{a \in \mathcal{A}} a^*$

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a point x is in $\partial U_{\mathcal{A}}$ if $x \subset \mathcal{A}^+$ of course $\mathcal{A}^+ = \{X \subset \mathbb{N} : \{a \in \mathcal{A} : X \cap a \neq^* \emptyset \text{ is infinite}\}\}$ Say that an open $U \subset \mathbb{N}^*$ is an adf* if there is an infinite adf \mathcal{A} such that $U = U_{\mathcal{A}} = \bigcup_{a \in \mathcal{A}} a^*$ i.e. U is paracompact (and not compact)

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 $\mathcal A$ is completely separable if each $X\in\mathcal A^+$ contains* some $a\in\mathcal A$

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Simon: tfae

- 1. There is no maximal nwd subset of \mathbb{N}^\ast
- 2. every madf has a completely separable madf refinement

more adf terminology

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more adf terminology

Definition

Let $\mathcal{A}_0, \ldots, \mathcal{A}_n$ be adf's

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more adf terminology

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• for $X \in \mathcal{A}_0$, $\mathcal{A}_0 \upharpoonright X = \{a \cap X : a \in \mathcal{A}_0\} \setminus [\mathbb{N}]^{<\aleph_0}$

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Let $\mathcal{A}_0, \ldots, \mathcal{A}_n$ be adf's • for $X \in \mathcal{A}_0, \mathcal{A}_0 \upharpoonright X = \{a \cap X : a \in \mathcal{A}_0\} \setminus [\mathbb{N}]^{<\aleph_0}$ • $\mathcal{A}_1 \prec \mathcal{A}_0$ if $\mathcal{A}_1 = \bigcup \{\mathcal{A}_1(a) = \mathcal{A}_1 \cap [a]^{\aleph_0} : a \in \mathcal{A}_0\}$ • $\mathcal{A}_1 \prec^+ \mathcal{A}_0$ if also $\mathcal{A}_0^+ \subset \bigcup \{\mathcal{A}_1(a)^+ : a \in \mathcal{A}_0\}$ (corresponds to $\partial U_{\mathcal{A}_0}$ is nwd in $\partial U_{\mathcal{A}_1}$) • $\mathcal{A}_1 \prec^{++} \mathcal{A}_0$ if also each $\mathcal{A}_1(a)$ is a madf on a

Let A₀,..., A_n be adf's
for X ∈ A₀, A₀ ↾ X = {a ∩ X : a ∈ A₀} \ [N]^{<ℵ₀}
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A₁ ≺⁺ A₀ if also A⁺₀ ⊂ ∪{A₁(a)⁺ : a ∈ A₀} (corresponds to ∂U_{A₀} is nwd in ∂U_{A₁})
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every infinite completely separable adf \mathcal{A} is +-partitionable and +-refinable because $\mathfrak{c} = |\{a \subset^* X : a \in \mathcal{A}\}|$ for $X \in \mathcal{A}^+$

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Some difficulties

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If for <u>all</u> madf A₀ there is A₁ ≺⁺ A₀, then there is a completely separable madf (which is presently unknown)

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Questions

• Do there exist madf's with $A_2 \prec^+ A_1 \prec^+ A_0$?

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Questions

- Do there exist madf's with $A_2 \prec^+ A_1 \prec^+ A_0$?
- **2** Can A_1 also be +-partitionable?

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If $\mathfrak{a} = \aleph_1$ (or $\mathfrak{a} = \mathfrak{h} = cof([\mathfrak{h}]^{\aleph_0})$), then $\mathcal{A}_1 \prec^+ \mathcal{A}_0$ exists. but unlikely that $|\mathcal{A}_1| = \mathfrak{a}$ so no continuing

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Proof.

If $\mathcal{A}_0 = \{a_\alpha : \alpha \in \omega_1\}$, then for each α choose an almost disjoint refinement \mathcal{X}_α for $(\{a_\beta : \beta < \alpha\})^+$

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Main Lemma

Assume we have adf's $\{\mathcal{A}_t: t\in m^{\leqslant n}\}$ satisfying for $t\in m^{\leqslant n}$

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$$\{ \mathcal{A}_{t \frown i} : i < m \} \text{ is a } + \text{-partition of } \bigcup_{i < m} \mathcal{A}_{t \frown i} \\ \text{ hence } \mathcal{A}_{t \frown i} \prec^+ \mathcal{A}_t$$

e.g.
$$\bigcup \{ \mathcal{A}_t : t \in m^k \}$$
 is a madf for each $k \leq n$

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Corollary

If there is a completely separable madf, then $\{A_t : t \in m^{\leq n}\}$ as above exists for all n, m,

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Corollary

If there is a completely separable madf, then $\{A_t : t \in m^{\leq n}\}$ as above exists for all n, m, hence \mathbb{N}^* will map onto every $m^{\leq n}$ by an open continuous map.

Question

For which n, m does such a family $\{A_t : t \in m^{\leq n}\}$ exist? Are there natural ZFC constructions? Is this equivalent to the existence of a completely separable madf?

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Definition

Let
$$T_{n,m} = m^{\leqslant n} \cup \{t^\frown m : t \in m^{\leqslant n}\} \subset (m+1)^{\leqslant n}$$

i.e. the subtree of $(m+1)^{\leqslant n}$ such that having m in the range makes it a maximal node.





Image: A matrix and a matrix

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There is an infinite completely separable adf A.

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There is an infinite completely separable adf A.

Lemma

There exist adf's $\{A_t : t \in m^{\leq n}\}$ satisfying for $t \in m^{< n}$

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There is an infinite completely separable adf A.

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$$\bigcirc \bigcup_{i < m} \mathcal{A}_{t \frown i} \prec^{+} \mathcal{A}_{t} \ (a \in \mathcal{A}_{t} \ \mathsf{NOT} \text{ refined by a madf})$$

$$\{ \mathcal{A}_{t \frown i} : i < m \} \text{ is a } + \text{-partition of } \bigcup_{i < m} \mathcal{A}_{t \frown i} \\ \text{ still have } \mathcal{A}_{t \frown i} \prec^{+} \mathcal{A}_{t}$$

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Proof.

Same construction except that, for $t \in m^{\leq n}$ and $a \in A_t$, $\bigcup_{i \leq m} A_{t^{\frown}i}(a)$ is a completely separable adf but not mad
LAST SLIDE!!!

Theorem

There is an open continuous map from \mathbb{N}^* onto $T_{n,m}$

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Proof.

Start with $\{\mathcal{A}_t : t \in (m+1)^{\leq n}\}$ and for $t \in m^{\leq n}$, $U_t = U_{\mathcal{A}_t} \text{ AND } U_{\mathcal{A}_t \frown m} \subset U_{t \frown m} = \bigcup_{a \in \mathcal{A}_t} a^* \setminus \operatorname{cl} (\bigcup_{i < m} U_{t \frown i})$

Loosely speaking: $U_{t \frown m}$ absorbs the missing non-madness part of each $a \in A_t$ and makes up for the fact that we weren't using a completely separable madf at each step.

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Now for some finite topology!

 $T_{n,m}$ maps onto $m^{\leq n}$ by an open continuous map. Solving the Modal Logic problem.